

# Embedding Intensional Semantics into Inquisitive Semantics

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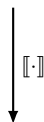
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Semantic model: object language with model-theoretic denotation

- Object language: simply typed  $\lambda$ -calculus
- Denotation: sets and functions
- Map: **compositional interpretation**

Object language



Denotation

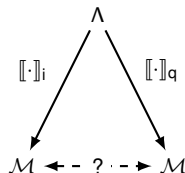
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We investigate the relation between:

**Intensional interpretation**  $[[\cdot]]_i$  (declarative)

vs. **Inquisitive interpretation**  $[[\cdot]]_q$  (declarative + interrogative)

- 1 Intensional and inquisitive semantics
  - Object language
  - Intensional interpretation
  - Inquisitive interpretation
- 2 Inquisitivation
- 3 Object language modification

# Object language

## Simply typed $\lambda$ -calculus

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- linguistic constants
  - Mary: **m** : IND, John: **j** : IND, Ash: **a** : IND
  - **sleep** : IND  $\rightarrow$  PROP
  - **try** : IND  $\rightarrow$  (IND  $\rightarrow$  PROP)  $\rightarrow$  PROP
- logical connectives
  - $\neg$  : PROP  $\rightarrow$  PROP
  - $\wedge$  : PROP  $\rightarrow$  PROP  $\rightarrow$  PROP
  - $\vee$  : PROP  $\rightarrow$  PROP  $\rightarrow$  PROP
  - ...

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### ■ Example:

(1) Mary sleeps.

**sleep m** : PROP

(2) Mary does not sleep.

$\neg$  (**sleep m**) : PROP

(3) Mary tries to sleep.

**try m** ( $\lambda x$ . **sleep x**) : PROP



# Intensional interpretation

## Intensional model:

- truth values  $\{0, 1\}$
- individuals  $D = \{m, j, a\}$
- possible world  $W = \{M, J, A\}$

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$$[[\mathbf{sleep} \mathbf{m}]]_i = [[\mathbf{sleep}]]_i ([[ \mathbf{m} ]]) = \{M\}$$

$$[[\neg(\mathbf{sleep} \mathbf{m})]]_i = W \setminus [[\mathbf{sleep} \mathbf{m}]]_i = \{J, A\}$$

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Hamblin's **alternative semantics** [4]:

- interrogative proposition = **set of answers** (i.e. intentional proposition), e.g.

$$\begin{aligned} \llbracket \text{Does Mary sleeps?} \rrbracket_q &= \{ \llbracket \text{Mary sleeps.} \rrbracket_i, \llbracket \text{Mary does not sleep.} \rrbracket_i \} \\ &= \{ \{M\}, \{J, A\} \} \end{aligned}$$

- several drawbacks [2]



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Representing a question by the **downward-closed** set of its answers:

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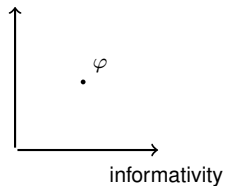
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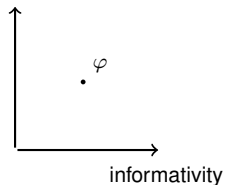
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- purely informative issue : only 1 alternative, e.g.

$$\llbracket \text{Mary sleeps.} \rrbracket_q = \{ \{M\} \}$$

→ uniform account of declaratives and interrogatives

# Question

Suppose we have a complete intensional lexicon

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## Practical need

Can we design an inquisitive interpretation **out of** the intensional one?

# Inquisitive interpretation

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$$\begin{aligned} &\text{with } P : D \rightarrow \mathcal{P}(\mathcal{P}(W)) \\ \text{but } \llbracket \mathbf{try} \rrbracket_i &: D \rightarrow (D \rightarrow \mathcal{P}(W)) \rightarrow \mathcal{P}(W) \end{aligned}$$

# Goal

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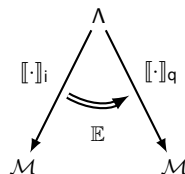
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## Inquisitivation:

- Embedding intensional semantics into inquisitive semantics



1 Intensional and inquisitive semantics

2 Inquisitivation

- Inquisitivation
- Properties

3 Object language modification

# First ideas

$\mathbb{E}$  indexed by types of  $\Lambda$ , following [3]

$$\mathbb{E}_{\text{IND}} : D \rightarrow D$$

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On simple examples:

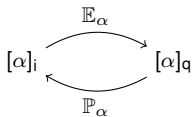
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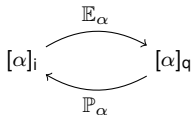
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Using

$$\begin{array}{ccc}
 \mathbb{U} : & \mathcal{P}(\mathcal{P}(W)) & \rightarrow & \mathcal{P}(W) \\
 & \mathcal{I} & \mapsto & \bigcup_{p \in \mathcal{I}} p
 \end{array}$$

Property:

$$\mathbb{U}(\mathcal{P}(p)) = p$$

## Definition of inquisitivation

Complete definition:

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## Main theorem

Let  $\varphi$  be a proposition. Then,  $\models_q \varphi$  if and only if  $\models_i \varphi$ .

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# Inquisitive logic

Formulas:

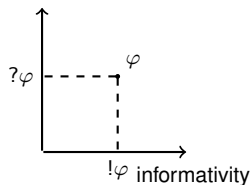
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inquisitiveness



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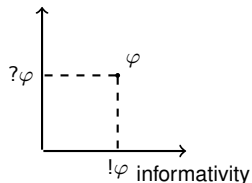
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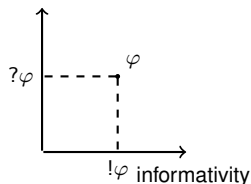
$$\llbracket \varphi \vee \psi \rrbracket = \llbracket \varphi \rrbracket \cup \llbracket \psi \rrbracket$$

(5) Does Mary sleep?

$$?(\mathbf{sleep\ m}) = (\mathbf{sleep\ m}) \vee (\neg(\mathbf{sleep\ m}))$$

$$\llbracket ?(\mathbf{sleep\ m}) \rrbracket = \{\{M\}, \emptyset\} \cup \{\{J, A\}, \{J\}, \{A\}, \emptyset\}$$

inquisitiveness



## What about logical constants?

Strong inquisitive interpretation  $\llbracket \cdot \rrbracket_{sq}$

- same, but logical constants interpreted as in  $Inq$

## What about logical constants?

### Strong inquisitive interpretation $\llbracket \cdot \rrbracket_{sq}$

- same, but logical constants interpreted as in Inq
- $\llbracket \wedge \rrbracket_q \neq \llbracket \wedge \rrbracket_{sq}$
- but if  $\mathcal{I}, \mathcal{Q}$  purely informative

$$\llbracket \wedge \rrbracket_q (\mathcal{I})(\mathcal{Q}) = \llbracket \wedge \rrbracket_{sq} (\mathcal{I})(\mathcal{Q})$$

- $\llbracket \wedge \rrbracket_{sq}$  better generalization of  $\llbracket \wedge \rrbracket_i$



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### Strong inquisitive interpretation $\llbracket \cdot \rrbracket_{sq}$

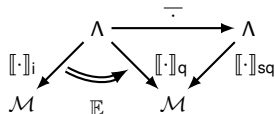
- same, but logical constants interpreted as in  $Inq$
- $\llbracket \wedge \rrbracket_q \neq \llbracket \wedge \rrbracket_{sq}$
- but if  $\mathcal{I}, \mathcal{Q}$  purely informative

$$\llbracket \wedge \rrbracket_q (\mathcal{I})(\mathcal{Q}) = \llbracket \wedge \rrbracket_{sq} (\mathcal{I})(\mathcal{Q})$$

- $\llbracket \wedge \rrbracket_{sq}$  better generalization of  $\llbracket \wedge \rrbracket_i$

Looking for: **object language translation**  $\bar{\cdot}$  such that

- $\llbracket \varphi \rrbracket_q = \llbracket \bar{\varphi} \rrbracket_{sq}$ , for every  $\varphi$



## What about logical constants?

### Strong inquisitive interpretation $\llbracket \cdot \rrbracket_{sq}$

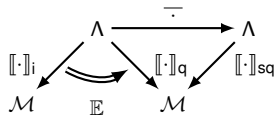
- same, but logical constants interpreted as in Inq
- $\llbracket \wedge \rrbracket_q \neq \llbracket \wedge \rrbracket_{sq}$
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$$\llbracket \wedge \rrbracket_q (\mathcal{I})(\mathcal{Q}) = \llbracket \wedge \rrbracket_{sq} (\mathcal{I})(\mathcal{Q})$$

- $\llbracket \wedge \rrbracket_{sq}$  better generalization of  $\llbracket \wedge \rrbracket_i$

Looking for: **object language translation**  $\overline{\cdot}$  such that

- $\llbracket \varphi \rrbracket_q = \llbracket \overline{\varphi} \rrbracket_{sq}$ , for every  $\varphi$



Syntactic translation

$$\overline{\varphi \vee \psi} =!(\overline{\varphi} \vee \overline{\psi})$$

$$\overline{\exists x. \varphi} =!(\exists x. \overline{\varphi})$$

$$\overline{c} = c \quad \text{for the other constants}$$

$$\overline{tu} = \overline{t} \overline{u}$$

$$\overline{\lambda x. t} = \lambda x. \overline{t}$$

# Conclusion

- **Conservative extension** from intensional semantics to inquisitive semantics

## Future prospects

- Using inquisitivation to define higher-order inquisitive logic

- [1] I. Ciardelli, J. Groenendijk, and F. Roelofsen. Inquisitive Semantics: A New Notion of Meaning. **Language and Linguistics Compass**, 7(9):459–476, 2013. ISSN 1749-818X. doi: 10.1111/lnc3.12037.
- [2] I. Ciardelli, F. Roelofsen, and N. Theiler. Composing alternatives. **Linguistics and Philosophy**, 40(1):1–36, February 2017. ISSN 1573-0549. doi: 10.1007/s10988-016-9195-2.
- [3] Ph. de Groote and M. Kanazawa. A Note on Intensionalization. **Journal of Logic, Language and Information**, 22(2):173–194, April 2013. ISSN 1572-9583. doi: 10.1007/s10849-013-9173-9.
- [4] Ch. L. Hamblin. Questions in Montague English. **Foundations of Language**, 10(1):41–53, 1973. ISSN 0015-900X.

Interpretation of  $\lambda$ -terms:

$$\llbracket x^\alpha \rrbracket_{i, \xi} = \xi_\alpha(x)$$

$$\llbracket c^\alpha \rrbracket_{i, \xi} = \mathcal{I}_\alpha(c)$$

$$\llbracket t^{\alpha \rightarrow \beta} u^\alpha \rrbracket_{i, \xi} = \llbracket t^{\alpha \rightarrow \beta} \rrbracket_{i, \xi} (\llbracket u^\alpha \rrbracket_{i, \xi})$$

$$\llbracket \lambda x^\alpha. t^\beta \rrbracket_{i, \xi} = \mathbf{a} \in [\alpha]_i \mapsto \llbracket t^\beta \rrbracket_{i, \xi[x^\alpha := \mathbf{a}]}$$